

a) $L_4 = \{ a, a^2, a^3, a^4, b^2a, b^2a^2, ab^2a \}$ 7 x 1/2 = 3.5 pts

b) $D_1 = a D_2 + b D_4$

$D_2 = a D_5 + b D_4 + \epsilon$

$D_3 = a D_2 + b D_4$

$D_4 = a D_6 + b D_3$

$D_5 = a D_5 + b D_4 + \epsilon$

$D_6 = (a+b) D_6$

6 x 1/2 = 3 pts

1/2 x 3 = 1.5 pt

c) $D_6 = \phi$; $D_1 = D_3$ et $D_2 = D_5$

d) $D_5 = a^+ (b D_4 + \epsilon)$ 1 pt

$D_4 = b D_3$ 1 pt

$D_3 = b^2 D_3 + a D_2$

$D_3 = (b^2)^+ a D_2$ 1 pt

$D_2 = a^+ (b D_4 + \epsilon) + b D_4 + \epsilon$

$= (a^+ + \epsilon) b D_4 + a^+ + \epsilon$

$= a^+ b D_4 + a^+$

$D_2 = a^+ b b (b^2)^+ a D_2 + a^+$

$= a^+ (b^2)^+ a D_2 + a^+$

$D_2 = (a^+ (b^2)^+ a)^+ a^+$ 2 pts

$D_1 = a (a^+ (b^2)^+ a)^+ a^+$
 $+ (b^2)^+ a (a^+ (b^2)^+ a)^+ a^+$ 2 pts

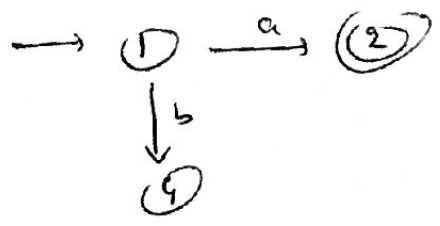
$L_2 (b^2)^+ a (a^+ (b^2)^+ a)^+ a^+$

2 pts

d) 9 pts

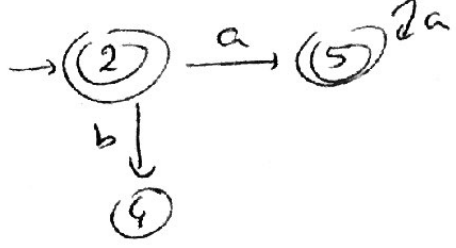
e) D_1 est le langage de l'automate

(2)



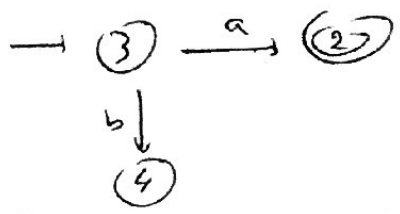
(0.5)

D_2 est le langage de l'automate



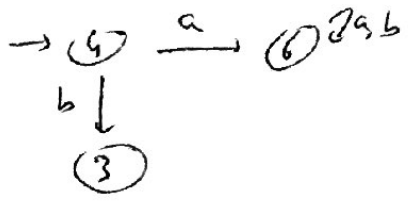
(0.5)

D_3 est le langage de l'automate



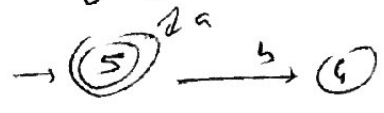
(0.5)

D_4 est le langage de l'automate



(0.5)

D_5 est le langage de l'automate



(0.5)

D_6 est le langage de l'automate $\rightarrow 6$ (0.5)

Les mots les plus courts de les langages sont respectivement:

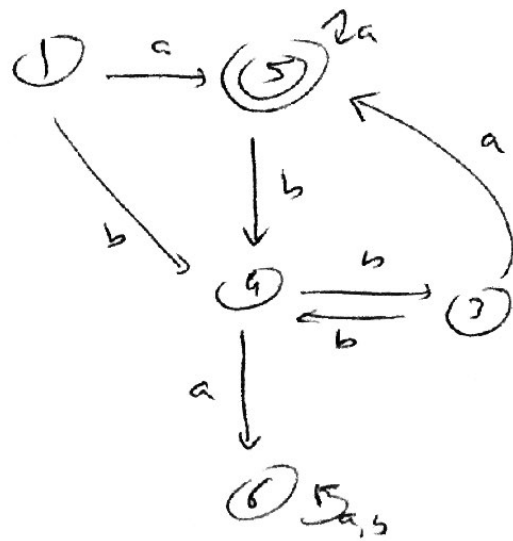
$m_1 = a; m_2 = \epsilon; m_3 = a; m_4 = \phi; m_5 = \epsilon; m_6 = \phi.$

(6 x 0.5 = 3 pts)

on a d'après c) $D_1 = D_3$ et $D_2 = D_5$ donc $q_1 \sim q_3$ et $q_2 \sim q_5$.

(0.5) + (0.5)

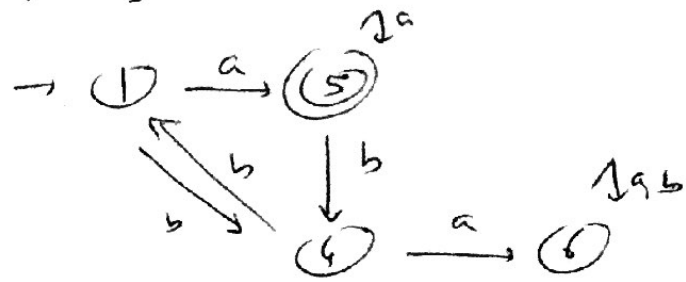
$q_2 \sim q_5$ donne!



(3)

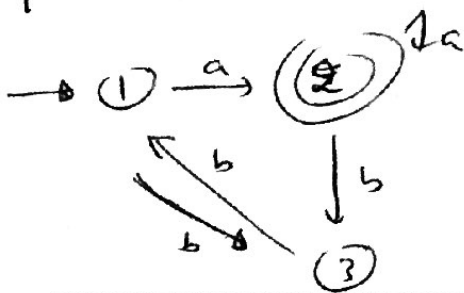
3pts

$q_1 \sim q_3$ donne



3pts

finalemment, l'automate minimal est dans sa forme incomplète réduite:



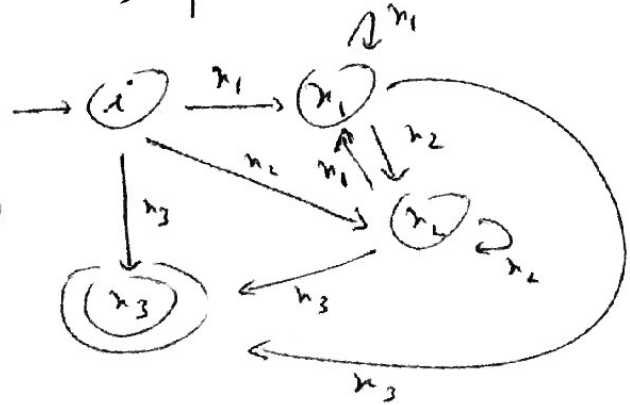
EX2 1a) $P = \{x_1, x_2, x_3\}$ (2pts)

$(a+b)^*c \equiv (x_1+x_2)^*x_3$ (1pt)

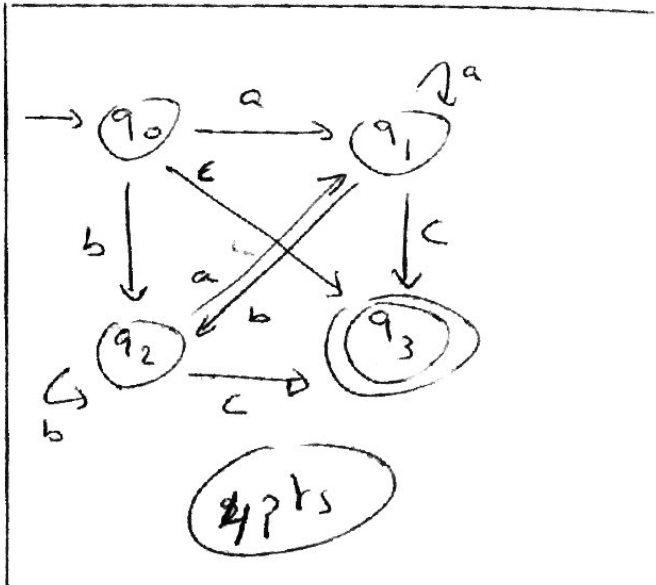
$D = \{x_3\}$ (2pts)

	Sub
x_1	x_1, x_2, x_3
x_2	x_1, x_2, x_3
x_3	-

1.5 + 1.5 = 3pts



3pts



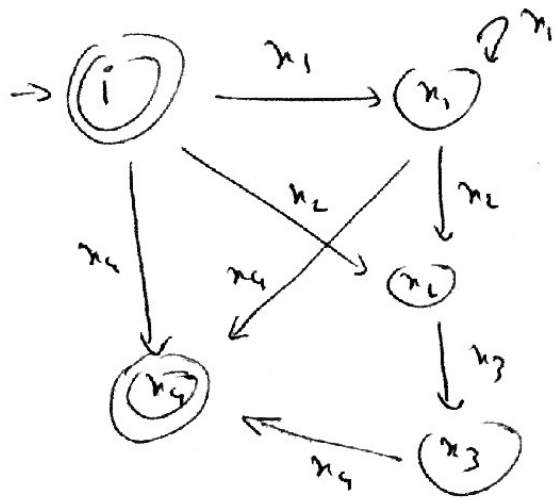
4pts

b) $\hat{a}^* (\epsilon + bb) a + \epsilon \equiv \pi_1^* (\epsilon + \pi_2 \pi_3) \pi_4 + \epsilon$

1pt (4)

$\text{Prz } \{ \epsilon, \pi_1, \pi_2, \pi_4 \}$ (2pts)

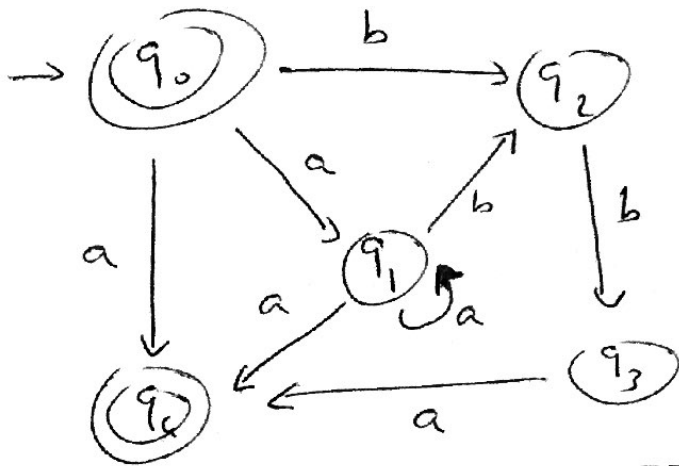
$\text{Drz } \{ \pi_4 \}$ (2pts)



3pts

	Sub
π_1	π_1, π_2, π_4
π_2	π_3
π_3	π_4
π_4	-

1.5
0.75 = 3pts
0.75



4pts

Ex 3 a) $S \rightarrow a s b^2 \rightarrow a a s b^2 b^2 \dots \rightarrow a^n s (b^2)^n$
 $\rightarrow a^n (b^2)^n$ et comme $S \rightarrow \epsilon$ aussi

$L = \{ a^n (b^2)^n / n \geq 0 \}$ (5pts)

b) $S \rightarrow a u c \rightarrow a a u b c \rightarrow a^2 a u b b c$
 $\rightarrow a a^2 u b^2 c \dots \rightarrow a a^n u b^n c$

$L = \{ a^{n+1} b^n c / n \geq 0 \}$ (5pts)

Ex 4.

a) $b a^2 = \epsilon b a^2 \epsilon \in a^* b^* a^* b^*$

(1 pt)

(5)

b) $b^* a^* \cap a^* b^* = \{ \epsilon \}$; $a^* \cup b^* = \{ a^n, b^m / n, m \geq 0 \}$ (1 pt)

$b^* a^* = \{ b^n a^m / n, m \geq 0 \}$

$= \{ b^n / n \geq 0 \} \cup \{ a^n / n \geq 0 \} \cup \{ b^n a^m / n, m \geq 1 \}$ (1 pt)

$a^* b^* = \{ a^n / n \geq 0 \} \cup \{ b^n / n \geq 0 \} \cup \{ a^n b^m / n, m \geq 1 \}$. (1 pt)

D'où $b^* a^* \cap a^* b^* = \{ a^n, b^m / n \geq 0, m \geq 0 \} = a^* \cup b^*$ (1 pt)

c) $a^* b^* \cap c^* d^* = \{ \epsilon \} \neq \emptyset$ (1 pt)

d) abcd est un mot de longueur 4 or les seuls mots de

longueur 4 que contient $(a(cd)^*b)^*$ sont:

acdb et abab

D'où $abcd \notin (a(cd)^*b)^*$.

(2 pts)

Ex 4: e) $u \in u^* \quad \text{dmc} (u+v)^* \subset (u^*+v)^*$ (1 pt)

reciproquement: $u^* \subset (u+v)^*$

$v \subset (u+v)^*$

$u^*+v \subset (u+v)^*$ (1 pt)

$(u^*+v)^* \subset ((u+v)^*)^* = (u+v)^*$ (1 pt)

D'où $(u+v)^* = (u^*+v)^*$ (0.5)

supplément
+ b) Soit $L = u+v$.

ona $L^* \subset u^* L^*$ car u^* contient ϵ .

enlevé de l'exam.

reciproquement: $u^* \subset (u+v)^* = L^*$

• donc $\mu^* L^* \subseteq L^* L^* = (L^*)^2 \subseteq L^*$.

(6)

D'où $(\mu + \nu)^* \subseteq \mu^* (\mu + \nu)^*$.

c) $\mu + \nu \subseteq \mu + \nu \mu^*$ donc $(\mu + \nu)^* \subseteq (\mu + \nu \mu^*)^*$

reciproquement:

$$\left\{ \begin{array}{l} \mu \subseteq (\mu + \nu)^* \\ \nu \mu^* \subseteq (\mu + \nu)^* \end{array} \right. \Rightarrow \mu + \nu \mu^* \subseteq (\mu + \nu)^*$$

donc $(\mu + \nu \mu^*)^* \subseteq ((\mu + \nu)^*)^* = (\mu + \nu)^*$

D'où $(\mu + \nu)^* = (\mu + \nu \mu^*)^* \rightarrow$ enlevé de l'exam

Ex4

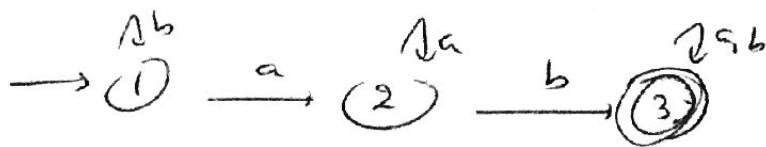
$$\text{f) } \left\{ \begin{array}{l} \mu \subseteq \mu^* \nu^* \\ \nu \subseteq \mu^* \nu^* \end{array} \right. \Rightarrow (\mu + \nu) \subseteq \mu^* \nu^* \Rightarrow (\mu + \nu)^* \subseteq (\mu^* \nu^*)^*$$

aussi $\mu^* \nu^* \subseteq (\mu + \nu)^*$ donc $(\mu^* \nu^*)^* \subseteq (\mu + \nu)^{**}$

$$(\mu^* \nu^*)^* \subseteq (\mu + \nu)^*$$

D'où $(\mu + \nu)^* = (\mu^* \nu^*)^*$ (0.5)

Ex5 soit l'automate A:



(10 pts)

• sous la forme complète; son complémentaire



(10 pts)

$$\bar{L} = \epsilon + b^* a^+ + b^* = b^* a^*$$

(5 pts)